Outline and Shape Reconstruction in 2D ECCV 2022 TUTORIAL

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## Tutorial Outline

Intro \& Proximity Graphs
Curve Reconstruction
Benchmark \& Demo
Sketch Reconstruction
Visual Perception of Shapes
Shape Characterization

Stefan Ohrhallinger - 25 minutes
Stefan Ohrhallinger - 25 minutes, Q\&A 5 minutes
Amal Dev Parakkat - 25 minutes, break 15 minutes
Amal Dev Parakkat - 25 minutes, Q\&A 5 minutes
Jiju Peethambaran - 25 minutes
Jiju Peethambaran - 25 minutes, Q\&A 5 minutes

Topic: Intro \& Proximity Graphs

Motivation
Proximity Graphs

## Presenter:

## Stefan OHRHALLINGER

Researcher
Institute of Visual Computing \&
Human-Centered Technology

# Introduction 

Occl


The Problem


## Connect the Dots



Now try without the numbers
Reconstructed polygon

## Challenges for Curve Reconstruction



Non-uniform sampling


Noisy sampling


Sparse sampling


Outliers, multiple curves


Sharp Corners


Non-manifold curves

## How to Choose a Suitable Algorithm?



## A Benchmark Helps to Decide [OPP* ${ }_{21}$ ]



Evaluating algorithms on challenging curves, highlighting strengths \& weaknesses Quantitative analysis on: reconstruction quality \& run-time

## Scope of this Tutorial

We categorize 36 curve reconstruction algorithms:


Boundary samples


Area samples


Implicit curve


Polygonal curve

## Taxonomy of Algorithms



Graph-based


Non-manifold


Feature size based


HVS-based

Noisy fitting

Implicit


Sharp corners



Region Reconstruction

## Input Data: Properties



Non-uniform sampling: determines feature size


Noisy sampling: needs fitting

Outliers: needs filtering

## Reconstruction Output: Properties


$\operatorname{deg}(\mathrm{v})=2$
Manifold

$\operatorname{deg}(\mathrm{v}) \leq 2$
Open curves

Sharp Corners
Guarantees
$O(n \log n)$
Multiply Connected


Time Complexity

Input capabilities: e.g., noise, outliers, non-uniformity


Output capabilities: e.g., manifold, sharp, $O(n \log n)$
[Leeooa]
Connect2D
Peel $\quad$ Crawl

Concorde

$$
\alpha \text {-shapes }
$$

[Leno6]
FitConnect
StretchDenoise
$\gamma$-neighborhood
DISCUR
VICUR

> Shape-hull Graph
> ec-Shape
$\beta$-skeleton
Gathan
Conservative Crust
Ball-pivoting
[Aro98]

Robust HPR
Voronoi Labeling

Crust
[WYZ* ${ }_{14}$ ]
[AMoo]
NN-Crust

$$
\left[\mathrm{CFG}^{*} 05\right] \quad[\text { Rupi4 }]
$$

r-regular shapes
HNN-Crust
[FRor]
Optimal Transport [Gie99]

## EMST

[Hiyog] Edge exchanging

## Definitions Curve

Curve $\sum$ : Simple closed and planar
Smooth curve C: (collection of) twice-differentiable bounded i-manifolds $\in \mathbb{R}^{2}$
Sample set P: n points sampled on $\sum$ or $C$

## Definitions Sampling



## Definitions Sampling



Reach of a curve interval I: $\inf \operatorname{lfs}(p): p \in I \quad\left[O M W_{I} 6\right]$
$\rho$-Sampling [OMW ${ }_{16}$ ]:
$\forall \mathrm{p} \in \mathrm{C}, \exists \mathrm{s} \in \mathrm{S}:\|\mathrm{p}, \mathrm{s}\|<\rho$ reach $(\mathrm{p})$

## Proximity Graphs for a Point Set



GG



Minimum Spanning Tree: cycle-free graph spanning $P$ with minimum edge weights Relative Neighborhood Graph: $\forall(\mathrm{p}, \mathrm{q}): \mathrm{d}(\mathrm{p}, \mathrm{q}) \leq \mathrm{d}(\mathrm{p}, \mathrm{x}), \mathrm{d}(\mathrm{p}, \mathrm{q}) \leq \mathrm{d}(\mathrm{q}, \mathrm{x}) \forall \mathrm{x} \in \mathrm{P}, \mathrm{x} \neq \mathrm{p}, \mathrm{q}$ Gabriel Graph: All $(\mathrm{p}, \mathrm{q})$ with $\mathrm{p}, \mathrm{q} \in$ empty ball centered at $(\mathrm{p}, \mathrm{q})$

Delaunay Triangulation: circumcircles empty of P

## More Proximity Graphs



EMST ( $\mathrm{d} \geq \mathrm{I}$ ) $\rightarrow \mathrm{BC}_{\mathrm{o}}(\mathrm{d} \geq 2) \quad$ SIG edges: $\mathrm{r}=\left|\mathrm{v}, \mathrm{NN} \mathrm{N}_{\mathrm{I}}\right|$ overlap $\mathrm{DT} \backslash$ divergent concave

Boundary Complex
Connect2D [OMi3]

Sphere-of-Influence Graph
[Toussaint88]

Shape-Hull Graph
[PMI5]

Topic: Curve Reconstruction

Graph-based Algorithms
Feature size based Algorithms

## Presenter:

## Stefan OHRHALLINGER

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## Algorithms Based on Graphs - Overview

a-shapes [EKS83], Ball-pivoting [BB97]
$\beta$-skeleton [KR85]
y-neighborhood [Vel92]
Sculpting [Boi84a]
EMST [FMG94], edge exchange [OMir] and inflating [OMi3]
$r$-regular shape [Att97]
Shape-hull graph [PMisb], Voronoi labeling [PPT*is]
Crawl thru neighbors [PMi6]

## a-Shapes [EKS83]



# Disks of radius I/a 

Generalization of convex hull ( $\mathrm{a}=\mathrm{o}$ )

Extracting manifolds [BB97]
Later: Ball-pivoting algorithm [BMR*99]

## $\boldsymbol{\beta}$-Skeleton [KR85]


$\beta<I$

$\beta=2$ : Relative neighborhood graph
Empty lens formed with $\measuredangle \mathrm{prq}<\theta$
$=$ Intersection/Union of disks

$$
\theta= \begin{cases}\sin ^{-1} \frac{1}{\beta}, & \text { if } \beta \geq 1 \\ \pi-\sin ^{-1} \beta, & \text { if } \beta \leq 1\end{cases}
$$

## 8-Neighborhood [Vel92]

A unification of i2 graphs including convex hull, Delaunay triangulation, Gabriel graph, RNG, MST, nearest neighbor graph, a-shapes and $\beta$-skeletons.
$\gamma\left(\gamma_{0}, \gamma_{I}\right)$ is defined for $-I<\gamma_{0}, \gamma_{I}<I$ and $\left|\gamma_{0}\right| \leq\left|\gamma_{I}\right|$
Contains edges with empty neighborhood defined by disks using $\gamma_{0}, \gamma_{I}$

It can also reconstruct shapes not in the Delaunay graph

## Sculpting [Boi84a]

## EMST-based Reconstruction [FMG94]



Proves that EMST reconstructs (open) curve from sufficiently dense samples

## EMST-based Edge-Exchange Reconstruction [OMir]



Transform EMST with "snap" and "move" operations - combinatorial complexity

EMST-based Inflating Reconstruction [OMis]


EMST


BC: $\operatorname{deg}(\mathrm{v}) \geq 2$

$\operatorname{deg}(\mathrm{v}) \leq 2$
$\operatorname{deg}(\mathrm{v})=2$

## r-regular Shapes [Att97]

An $r$-regular shape has curvature $\geq r$ everywhere

Requires uniform sampling of boundary

Boundary consists of edges shared by Delaunay circumcircles with property of angle<threshold depending on uniform sampling density and curvature $r$

## Shape-Hull Graph [PMı5b]


(a). Divergent

(b). Non-divergent


Reconstructs smooth curves with divergent concavity
Eliminates Delaunay triangles with circumcenter outside boundary

## Voronoi Labeling [PPT*ig]



Incrementally labels orientation from estimated normals via Voronoi poles


Also computes medial axis

## Crawl Thru Neighbors [PMi6]



Connects neighbors greedily, heuristic decides curve closed/open
Parameter-free: handles open+multiple curves, holes and outliers

## Algorithms Based on Graphs - Conclusion

a-shapes [EKS83], Ball-pivoting [BB97], $\beta$-skeleton [KR85], $\boldsymbol{\gamma}$-neighborhood [Vel92], Sculpting [Boi84a], EMST [FMG94], edge exchange [OMir], inflating [OMi3], $r$-regular shape [Att97], Shape-hull graph [PMi5b], Voronoi labeling [PPT* ${ }^{\text {r9 }}$ ], Crawl thru neighbors [PMi6]

They often require a global parameter
Good results mostly for uniformly sampled point density
Delaunay graph is not guaranteed to contain the reconstruction
Reconstruction is often slow or trapped in local minima

## Algorithms Based on Feature Size - Overview

Crust [ABE98]
Anti-Crust [Gol99]
NN-Crust [DK99]
Conservative Crust [DMR99]
Lenz [Leno6]
Hiyoshi [Hiyog]
HNN-Crust [OMWI6]
SIGDT [MOW 22 ]

## Crust [ABE98]

Seminal paper: feature sized reconstruction - no more uniform sampling required

$\varepsilon$-Sampling [ABE98]:
$\forall p \in C, \exists s \in S:$
$\|p, \mathrm{~s}\|<\varepsilon \mathrm{Ifs}(\mathrm{p})$
extracts DG and Voronoi graph
Proof: $\varepsilon<0.252 \cong \alpha>15$ I $^{\circ}$

## Anti-Crust [Gol99]

Extracts the Crust in a single step from the Delaunay graph


Also extracts the medial axis skeleton

## NN-Crust [DK99]

Simple and elegant improvement of Crust:


First, connects point to nearest neighbor
Then to nearest neighbor in half-space s.t. angle $>90^{\circ}$
Proof: $\varepsilon<\mathrm{I} / 3$, corresponding to $\alpha>{\mathrm{I} 4 \mathrm{I}^{\circ}}^{\circ}$

## Conservative Crust [DMR99]



Crust


NN-Crust


Conservative Crust

Filters edges from Gabriel graph
Robust to outliers
Collections of open/closed curves
But requires a parameter
Misses some sharp corners

## Lenz: Probe Reconstruction [Leno6]



Starts with a seed edge and connects edges with a probe shape
Requires an angle parameter
Permits self-intersections
Claims $\varepsilon<0.48$ but no proof

## Hiyoshi: TSP [Hiyog]



Adapts Traveling Salesman Problem to multiple connected curves
Transforms it into maximum-weight 2 -factor problem (solvable in P time)
Proof for: $\varepsilon<1 / 3, \mathrm{u}<$ I. 46 (relative uniformity of adjacent edge lengths)

## HNN-Crust [OMWI6]

Simple variant of NN-Crust, reducing angle from $90^{\circ}$ to $60^{\circ}$ :


## SIGDT [MOW22]


г) $\mathrm{SIGDT}=\mathrm{SIG} \cap \mathrm{DT}$

2) Enforce $d \geq 2$

3) Inflating creates a manifold boundary

4) Sculpting interpolates interior vertices: $\mathcal{\varepsilon}<0.5, \mathfrak{u}<2$

Sampling Conditions of Crust Algorithms


NN-Crust:
$\varepsilon<\mathrm{I} / 3$

HNN-Crust:
$\rho<0.9$

## Algorithms Based on Feature Size - Conclusion

Crust [ABE98], Anti-Crust [Gol99], NN-Crust [DK99], Conservative Crust [DMR99]

Lenz [Leno6], Hiyoshi [Hiyog], HNN-Crust [OMW 16 ], SIGDT [MOW 22 ]

Guarantees on sampling condition
Work well for non-noisy point sets

## Outline

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Topic: Benchmark and Demo
To be precise:

- What all our benchmark has?
- How to use our benchmark?
- What all we evaluated?
- What are our conclusions?


## Presenter:

Amal Dev PARAKKAT<br>Assistant Professor<br>Institut Polytechnique de Paris<br>amal.parakkat@telecom-paris.fr

## The Benchmark

Our benchmark contains: Algorithms, Dataset, Sampling tools, Evaluation criterias, and Test scripts


## The Benchmark - Algorithms

## We included is publicly available algorithms

Contains algorithms from late gos (Crust family) to 2018

- Crust, NNCrust, CCrust, Gathan, GathanG, Lenz, Discur, Vicur, OptimalTransport, ConnectzD, Crawl, HNNCrust, FitConnect, StretchDenoise, Peel

We removed OptimalTransport from experiments since it simplifies curves


## The Benchmark - Dataset

Our dataset contains more than 2500 point sets:

- Classic data - Collected from various papers (using WebPlotDigitizer)
- Image data - Samples obtained from the silhouette images (taken from MPEG7 CE Shape-I, Edinburgh, Io7o-shape image databases)
- Synthetic data - Analytical (shapes with sharp corners, \& self-intersections) and $\varepsilon$-sampled points

Classic Data

|  |  |
| :---: | :---: |

Image Data


LFS sampling


Non-manifold


Sharp corners


## The Benchmark - Dataset

We also provide ground truths (linear approximation) as:

- Ordered vertices: A loop of vertices - for simple closed curves
- Edge list: List of edges - for complex curves

Grouped under the following categories:


Moreover, we provide an interactive ground truth generation tool

## The Benchmark - Sampling tools

LFS-sampling tool:

- Samples are made from input Bezier curve representation
- Maximal empty disks are computed to create a medial axis approximation
- Estimate LFS at each sample and use it to pick a set of samples satisfying the $\varepsilon$-sampling condition

Contour sampling tool:


A Bezier curve and its LFS-based sampling


A binary image
Extracted contour


## The Benchmark - Evaluation criteria

Let closest point correspondences be D and D' of two curves C and C'


$$
D=(s, t) \mid s \in C^{\prime}, t=M(s)
$$

$$
D^{\prime}=\left(s^{\prime}, t^{\prime}\right) \mid s^{\prime} \in C, t^{\prime}=M^{\prime}\left(s^{\prime}\right)
$$

where $M$ and $M^{\prime}$ be the respective non-bijective shortest distance maps
We use the following metrics to compare two curves:

$$
H_{D}\left(C, C^{\prime}\right)=\max \left\{\max _{(s, t) \in D}\|s-t\|, \max _{\left(s^{\prime}, t^{\prime}\right) \in D^{\prime}}\left\|s^{\prime}-t^{\prime}\right\|\right\}
$$

$$
\begin{aligned}
R M S_{D}\left(C, C^{\prime}\right)= & \sqrt{\frac{1}{N}\left(\sum_{(s, t) \in D}\|s-t\|^{2}+\sum_{\left(s^{\prime}, t^{\prime}\right) \in D^{\prime}}\left\|s^{\prime}-t^{\prime}\right\|^{2}\right)} \\
& \text { Root mean squared distance } \quad N=|D|+\left|D^{\prime}\right|
\end{aligned}
$$

## The Benchmark - Test scripts

Driver program can be run with various arguments and options
A set of test scripts for quantitatively \& qualitatively evaluate the algorithms
Each test script has a list of algorithms and test data, designed for the specific experiment:

- run-sampling.sh: $\varepsilon$-sampled [ABE98] test data
- run-noisy.sh: perturbed with uniform noise
- run-lfsnoise.sh: perturbed with lfs-based noise
- run-outliers.sh: added outlier points
- run-manifold.sh: whether reconstruction is a manifold
- run-sharp-corners.sh: sharp feature curves
- run-open-curves.sh: open curves
- run-multiple-curves.sh: multiply connected curves
- run-intersecting.sh: curves with intersections

Evaluation - Sampling density as $\varepsilon$-sampling


## Evaluation - Noise robustness (BB Diagonal)



## Evaluation - Noise robustness (LFS)



Evaluation - Noise robustness ( $\varepsilon$-sampling + LFS)

## 0.1

0.2
0.4

Ifs-varying Noisy sampling


## Evaluation - Outlier



## Evaluation - Running time



## Evaluation - Qualitative Comparison



## Evaluation - Qualitative Comparison



## Evaluation - Qualitative Comparison



## Evaluation - Qualitative Comparison



## Evaluation - Qualitative Comparison



## Evaluation - Qualitative Comparison



## Evaluation - Overview

In short, we evaluate the robustness of various algorithms based on:

- Sampling density as $\varepsilon$-sampling
- Noise robustness as $\delta$ of bounding box diagonal
- Noise robustness as $\delta$ of lfs
- Noise+sampling density as $\varepsilon$-sampling and $\delta$ of 1 fs
- Outliers robustness in \% of samples
- Average runtimes (in s)


## Evaluation - Summary

| Curve/Input feature | Best two algorithms in order |
| :--- | :--- |
| Uniform Noise | DISCUR, VICUR |
| Non-uniform Noise | STRETCHDENOISE, CONNECT2D |
| Outliers | HNN-CRUST, CRUST |
| Non-uniform sampling | HNN-CRUST, PEEL |
| Runtime | NN-CRUST, GATHAN1 |
| Manifold curves | CONNECT2D, CRAWL |
| Non-manifold curves | CRUST, LENZ |
| Sharp features | GATHANG, CONNECT2D |
| Open curves | VICUR, HNN-CRUST |
| Multiple curves | PEEL, HNN-CRUST |

## Outline

Topic: Sketch Reconstruction
To be precise:

- Sketching
- Sketching and reconstruction
- Sketch completion
- Rough sketch simplification


## Presenter:

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## Sketching

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Sketching is an integral part of everyone's life
Sketching comes across us in various stages of our life:

- Kids playing with pencils and paints
- Intrinsic part of various academic curriculums
- Professional usage

MOREOVER, IT'S FUN!!!!!


Relevance of sketch processing in Computer Graphics


Mountain creation


Printable models from VR drawings


3d modeling \& editing


Expressive sketch-based animation


Sketching animation in VR


## Sketching and Reconstruction

Sketching and reconstruction problem are closely related Reconstructing from:

- Simple scanned sketch

© de Goes et al.


## Sketching and Reconstruction

Sketching and reconstruction problem are closely related
Reconstructing from:

- Simple scanned sketch
- Missing strokes


© de Goes et al.


## Sketching and Reconstruction

Sketching and reconstruction problem are closely related
Reconstructing from:

- Simple scanned sketch
- Missing strokes
- Noisy sketch

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## Sketching and Reconstruction

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Sketching and reconstruction problem are closely related
Reconstructing from:

- Simple scanned sketch
- Missing strokes
- Noisy sketch
- Dirty sketches

© de Goes et al.


## Sketch completion and Sketch simplification

We look at two important subproblems:
Sketch completion and sketch simplification
We won't have a detailed discussion, but a very quick and brief overview of:

- A.D. Parakkat, P. Memari, M.P. Cani, "Delaunay Painting: Perceptual Image Colouring from Raster Contours with Gaps" - Computer Graphics Forum 2022
- A.D. Parakkat, P. Madipally, H.H. Gowtham, M.P. Cani, "Interactive Flat Coloring of Minimalist Neat Sketches" - Eurographics 2020 (Short paper)
- A.D. Parakkat, M.P. Cani, K. Singh, "Color by Numbers: Interactive Structuring and Vectorization of Sketch Imagery" - ACM CHI 202I
- A.D. Parakkat, U.B. Pundarikaksha, R. Muthuganapathy, "A Delaunay triangulation based approach for cleaning rough sketches" - Computers \& Graphics 2018


## Sketch completion

The input is a set of disconnected sketch strokes and is asked to appropriately connect them (in other words, filling the gaps in a line art)


## Sketch completion - Sketch coloring

Using Flood-fill algorithm


Distinct artistic feeling:

- Le Grand Méchant Renard et autres contes...
- Ernest \& Celestine


## Sketch completion - Sketch Coloring



COLOR HINTS
DESIRED COLORING
Infer the unknown boundary!!!

## Sketch completion - Delaunay Painting

Assumption - Required boundary is present in the Delaunay Triangulation (Delaunay confirming) Problem boils down to connecting appropriate points in Delaunay Triangulation


## Delaunay Painting - Defining flow



## Delaunay Painting - Defining flow


$\operatorname{Flow}(u, v)=\max (f(X): \forall$ paths $X$ from $u$ to $v)$

$$
f(X)=\min (\operatorname{Weight}(u, v): \forall(u, v) \in X)
$$

## Delaunay Painting - Defining flow



## Delaunay Painting - Defining flow


$F$ low $(u, v)=\max (f(X): \forall$ paths $X$ from $u$ to $v)$
$f(X)=\min (\operatorname{Weight}(u, v): \forall(u, v) \in X)$

## Delaunay Painting - Creating a graph



Weighted dual graph
Create a color_strength for all vertices and initialize it as 'o'

## Delaunay Painting - Logic

Recursive color spreading
Color strength updation
Priority queue based exploration


## Delaunay Painting - Updation

Color starting from a triangle ' v '


## Delaunay Painting - Demo



## Delaunay Painting - Demo



## Delaunay Painting - We can do a lot more



Are we missing something?
Too much work if the sketch has additional information like shading Not aesthetically appealing

## Delaunay Painting - Additional functionalities



How to address them?
Too much work if the sketch has additional information like shading - Diffuse colors Not aesthetically appealing - Give finishing using an aesthetic curve completion

## Delaunay Painting - Color Diffusion

Presence of shading information (hatching) makes the coloring process time consuming Artists usually use a comparatively smaller brush size for shading/hatching

## Bipartite the regions into two - colored and uncolored Recursively update the bipartition



Delaunay Painting - Aesthetic curve completion
Splitting points - Edges having different colors on its associated triangles Sharp corner or smooth curve?
Decision based on tangent approximation


## Delaunay Painting - Sharp corner heuristics

Angle constraint - angle between tangents less than $\pi / 3$
Perpendicular constraint - intersection to edge distance is less than $2^{*} \|$ edge $\|$ Linearity constraint - pixels near the endpoints are linearly arranged


All three constraints are qualified $\rightarrow$ sharp corner, else $\rightarrow$ smooth curve

## Delaunay Painting - SIMVC curves

Perceptually pleasing contour
Scale Invariant Minimum Variation Curve - to form more circular arcs

$$
E_{S I M V C-E n t e m}=\frac{\left(\int d s\right)^{5}}{\|B-A\|^{2}} \int\left(\frac{d \kappa(s)}{d s}\right)^{2} d s
$$



## Delaunay Painting - SIMVC curves

Perceptually pleasing contour
Scale Invariant Minimum Variation Curve - to form more circular arcs

$$
E_{S I M V C-E n t e m}=\frac{\left(\int d s\right)^{5}}{\|B-A\|^{2}} \int\left(\frac{d \kappa(s)}{d s}\right)^{2} d s \text { How small the curve is? }
$$



## Delaunay Painting - SIMVC curves

Perceptually pleasing contour
Scale Invariant Minimum Variation Curve - to form more circular arcs

$$
E_{S I M V C-E n t e m}=\frac{\left(\int d s\right)^{5}}{\|B-A\|^{2}} \int\left(\frac{d \kappa(s)}{d s}\right)^{2} d s \text { How much curved? }
$$



## Delaunay Painting - More Results



## Delaunay Painting - Labelling medical images



## Sketch Simplification



## Sketch Simplification - Automatic triangle growing

Delaunay triangles inside regions will have a "fat triangle"

Delaunay triangles inside adjacent strokes will have only "thin triangles"
 Overview:


## Automatic triangle growing

Algorithm: Start from the largest "valid" ungrouped triangle
Recursively group neighboring triangles until a "condition" is satisfied Restart the procedure

(a)

(g)

(b)

(h)

(c)

(i)

(d)

(j)

(e)

(k)

(f)

(1)

## Automatic triangle growing



Stop the procedure when there are no more "valid" ungrouped triangles Pick all the ungrouped triangles (lies inside adjacent strokes) - group and color them Compute the skeleton of this colored group, and fit cubic Bezier curves

## Sketch Simplification - Color by numbers

Perception plays an important role in simplification - Not available in Automatic triangle grouping Design sketches usually have construction lines Idea: Make users annotate the parts that should be grouped

Playful interface: Ask the user to give same color on the opposite sides of a required stroke

And we already know how to do it!!! - Delaunay painting


## Color by numbers - Complete procedure



## Color by numbers - Demo



## Color by numbers - Results

Reconstruction from Sketches


## Outline

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Topic: Shape Characterization

- Shape/Region Reconstruction
- HVS based Algorithms
- Delaunay based Algorithms
- Sampling Models
- Evaluation Practices
- Future Directions

Presenter:

Jiju Peethambaran
Assistant Professor
Saint Mary's University, Halifax

## Shape Characterization or Region Reconstruction

- Given a finite set of points sampled from a planar object or region, construct a polygonal boundary that best approximates the object or region

- Inputs are known as dot pattern/area samples/region samples
- Outputs: graphs or polygons
- Compared to curve reconstruction, more signals (or samples) about the shape is available


## Applications

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- Computer graphics- geometric modeling ${ }^{[1]}$

(a)


(d)

Figure : (a) Points on surface with constraints, (b) Points in parametric 2D space, (c) Reconstruction, (d) Trimmed patch

- Identification of island failure regions in the design space of reliability-based crash optimization ${ }^{[2]}$

[1] Sundar et al. 2014, "Foot point distance as a measure of distance computation between curves and surfaces", Computers \& Graphics [2] Ganapathy et al. 2015, Alpha shape-based design space decomposition for island failure regions in reliability-based design", Struct. Multidisc.


## Applications

- Map generalization- e.g., aggregation of buildings to form single polygon ${ }^{[3]}$



## - Outline of trees or flock of birds ${ }^{[4]}$


[3] Roth et al. 2014, "a typology of operators for maintaining legible map designs at multiple scales", Cartographic Perspectives [4] Pandey et al. 2021, "Towards Video based Collective Motion Analysis through Shape Tracking and Matching", IET Electronic Letters

## Challenges

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- III-posed/Vague problem ${ }^{[5]}$ - a precise mathematical definition for 'shape' is almost impossible
— Rich variety of shapes and forms
— Heterogeneity of point set sampling (density and distribution)
- Different interpretations for 'shape' based on human cognition, visual perception and application demands

[5] Edelsbrunner 1998, "Shape Reconstruction with Delaunay Complex", LATIN


## Region Reconstruction Criteria ${ }^{[6]}$

- Should every member fall within the region or outliers permitted?

- Should any points fall on the boundary, or they must fall in the interior?

[6] Galton et al. 2006, "What is the Region Occupied by a Set of Points?", GIScience


## Region Reconstruction Criteria ${ }^{[6]}$

- Should the region boundary be polygonal, or can it be smooth and curved?

- Should the region boundary be a simple polygon?

[6] Galton et al. 2006, "What is the Region Occupied by a Set of Points?", G/Science


## Classification



## Human Visual Perception

- Gestalt Laws of visual perception

Proximity



## Closure



## DISCUR ${ }^{[7]}$

- A vision function that encodes p's relation to $T_{\mathrm{q}}$ and the edge $(r, q)$

$$
E\left[p, T_{q}\right]=h_{d} \frac{h}{s}\left(1+\frac{h_{d}}{\sigma_{d}}\right)^{\frac{\sigma_{d}}{h_{d}}}
$$

- If $\mathrm{d}(\mathrm{p}, \mathrm{q})<\mathrm{E}\left[\mathrm{p}, \mathrm{T}_{\mathrm{q}}\right]$, connect p to q
- Parameter free algorithm

- Open/closed curves, multiple curves, sharp corners
- Dense sampling at sharp corners



A failure case

## $\mathrm{VICUR}^{[8]}$

- DISCUR limitation: arbitrary selection of candidate points
- Vision function that encodes proximity and continuity

$$
E\left[p, T_{p_{1}}\right]=\left[c\left(\frac{\alpha_{\mathrm{s}}}{\bar{\alpha}}-1\right)^{2}+\left(\frac{1-c}{4}\right)\left(\frac{d_{\mathrm{s}}}{\bar{d}+\sigma}\right)^{2}+1\right]^{-1}
$$



- Candidate point with highest $E$ value is connected
- Sensitive to parameters, e.g., c balances the smoothness and nearness


VICUR result [8]

## Simple Shape ${ }^{[9]}$

- Shape according to HVS based on how the concavities are perceived
- Start from the convex hull and carve out the concavity by replacing outer edges by two new edges
- Edge selection is based on
— Closeness criteria
— Edge length criteria

— Angular constraints (angle (EAG)- angle(EGA) must be minimum)


## Simplicial Complex

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- k-simplex $\left({ }_{\sigma_{k}}\right)$ : non-degenerate convex hull of $\mathrm{k}+1$ geometrically distinct points in in $R^{d}$ where $k<=d$.

- Simplicial Complex:

A simplicial complex, $\mathcal{K}$ is a set containing finitely many simplices that satisfies the following two restrictions:

- $\mathcal{K}$ contains every face of every simplex in $\mathcal{K}$;
- For any two simplices, $\sigma, \tau \in \mathcal{K}$, their intersection $\sigma \cap \tau$ is either empty or a common face of $\sigma$ and $\tau$.


## Regular Simplicial Complex

## - Regular 2-simplicial complex:

A simplicial 2-complex $\mathcal{K}_{2}$ is said to be regular if it satisfies the following conditions:

- All the points in $\mathcal{K}_{2}$ are pairwise connected by a path on the edges.
- It does not contain any junction points, dangling edges or bridges.

(a)

(b)


## Delaunay Complex

- Given a finite set of points $S$ in $R^{d}$, Delaunay complex is a simplicial complex DT(S) consisting only of:

$\square$
all $d$-simplices whose circumspheres are empty of input points
— all $k$-simplices which are faces of other simplices in $\operatorname{DT}(S)$


## Alpha Shape

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- "Shape formed by a set of points"
- Ice Cream Carving Analogy ${ }^{[\$]}$
( Ice cream mass occupied in $\mathrm{R}^{\mathrm{d}}$ and chocolate points
- Sphere formed ice cream spoon

C Carve out ice cream without bumping into the chocolate points
$\square$ Carving spoon of small radius $\square$ points
$\square$ Carving spoon with huge radius $\square$ convex hull


Image courtesy: CGAL Alpha shapes
[\$] H. Edelsbrunner and E. P. Mücke. Three-dimensional alpha shapes. ACM Trans. Graph., 13(1):43-72, January 1994.

## Alpha Shape ${ }^{\text {Iol }}$

- $\alpha_{-}$exposed simplex: A k-simplex is $\alpha$ - exposed if there exists an empty

$$
\sigma_{k}=\partial b \cap S
$$



DEFINITION The boundary $\partial S_{\alpha}$ of the $\alpha$-shape of the point set $S$ consists of all $k$-simplices of S for $0 \leq k<d$ which are $\alpha$-exposed,

$$
\partial S_{\alpha}=\left\{\sigma_{k} \mid k \leq d,\left(v_{0}, v_{1}, \ldots, v_{k}\right) \subseteq S \text { and } \sigma_{k} \text { are } \alpha-\text { exposed }\right\}
$$


[10] Edelsbrunner et al. 1983, "On the shape of a set of points in the plane", IEEE Transactions on Information Theory
[11] Fischer K., "Introduction to Alpha Shapes", Technical Report, Stanford University

## A-Shape ${ }^{[12]}$

- Let $\mathcal{A} \subseteq R^{2}$ is a finite set of points, S
- $\mathcal{A}$-shape ${ }^{\text {is generated by connecting }} p, q \in S$ if there is an empty circle that ravoco inrough p, q and a

$$
\in \mathcal{A}
$$

- Two parameter family of point sets

$$
\mathcal{A}=\mathcal{A}(\alpha, t)
$$

$\square_{t \in[0,1]}$ is a local density measure $t \in[0,1]$
$\square_{\alpha \geq 0}$ level of detail of the shape


## Chi Shape ${ }^{\left[{ }^{2}\right]}$

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- Simple polygon that characterizes the shape of point set, S.
- Start with the Delaunay Triangulation of S.
- Repeatedly remove longest boundary edges greater than a threshold / subjected to regularity constraints.
- Generates a regular polygon that contains $S$.



## Chi Shape

- How to select I?

$$
\lambda_{P}= \begin{cases}1 & \text { if } l \geqslant \max _{P} \\ \frac{l-\min _{P}}{\max _{P}-\min _{P}} & \text { if } \min _{P} \leqslant l<\max _{P} \\ 0 & \text { if } l<\min _{P}\end{cases}
$$

- Good characterization via normalized length parameier <ip $0.27<0.29$

$0.77<\lambda_{P} \leqslant 1.00$

 half-way between max-MST and min-MAX?


## Characterization of object boundaries: Divergent Concavity ${ }^{1 / 4}$ =CCl <br> TEL AVIV 2022

- Closed, planar and positively oriented curve


## BT

- Inflection points and curvature
- Concave portion (green colored)
- BT-bi-tangent, BTP-bi-tangent points



## Characterization of object boundaries: Divergent Concavity

- Closed, planar and positively oriented curve

BT

- Inflection points and curvature
- Concave portion (green colored)
- BT-bi-tangent, BTP-bi-tangent points
- Pseudo-concavity



## Characterization of object boundaries: Divergent Concavity

- Extremal Vs Non-extremal BT


Characterization of object boundaries: Divergent Concavity


- Divergent pseudo-concavity


## Characterization of object boundaries: Divergent Concavity



- If all the pseudo-concavities are divergent, then the curve is divergent


## Divergent Concavity

- Implications ${ }^{[19]}$



Triangles in divergent concave region

## Characterization of object boundaries: Divergent Concavity

- Triangles in divergent concave regions are:
- Obtuse
- Longest edge facing towards the extremal BT



## Relaxed Gabriel Graph ${ }^{[5]}$

- Consists of all Gabriel edges and a few non-Gabriel edges
- RGG(S) retains a non-Gabriel edge ( $p, q$ ) of DT(S) if it satisfies either of the following:
- Circumcenter of the Delaunay triangle $\triangle p q r$ for which $(p, q)$ is the characteristic edge, lies internal to $\partial R G G(S)$.
- Removal of $(p, q)$ violates regularity in $\operatorname{RGG}(\mathrm{S})$.

[15] Peethambaran et al. 2015, "A nonparametric approach to shape reconstruction from planar point sets through Delaunay filtering", Computer Aided Design


## Relaxed Gabriel Graph

- Hole structure: fat triangle surrounded by sets of thin triangles

(d)

(g)
(h)
(i)


## Relaxed Gabriel Graph

- Order the boundary triangles based on their circum-radii (priority queue)
- Remove the boundary triangles if they are deletable
deletable $\square$ circum-center lie outside the intermediate boundary and the removal does not violate regularity of the simplicial complex.
- $O(n \log n)$ complexity



## EC-shape ${ }^{[16]}$

- Exterior triangle and exterior edge
- Circle constraint: Non-empty diametric, chc

[16] Methirumangalath et al. 2015, "A unified approach towards reconstruction of a planar point set.", Computers \& Graphics


## EC-shape

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- Remove the exterior edges if it satisfy circle constraints and regularity constraints
- Illustration:
- Construct Delaunay



## EC-shape

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- Remove the exterior edges if it satisfy circle constraints and regularity constraints
- Illustration:
- Non-empty diametric circle
- Empty diametric circle and non-empty midpoint circles


(f)

(g)


## EC-shape

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- Remove the exterior edges if it satisfy circle constraints and regularity constraints
- Illustration:
- Empty diametric and non-empty chord circles
- All circles empty

(k)


(1)


## EC-shape

- Final shape
- Under r-sampling, EC-shape is
homeomorphic to a simple closed curve



# CT-shape ${ }^{\left[{ }_{77}\right]}$ 

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- Coordinated triangles: If the circumcenters of neighboring triangles lie on the same half plane made by the shared edge
- Skinny triangles: non-obtuse triangle with ba
 le distance between its circumcenter and incenter
- Degree constraints: for vertices with c..............rter edges are retained.
[17] Thayyil et al. 2020, "An input-independent single pass algorithm for reconstruction from dot patterns and boundary samples.", Computers Aided Geometric Design


## CT-shape

- Mark all the shared edges of the coordinated triangles, two longer edges of the skinny triangles
- Create a graph consisting of all unmarked edges
- Apply degree constraints to get the final shape

- Theoretical guarantees under r-sampling


## Sampling Models: r-sampling

- A point set $S$ sampled from an object $O$ is said to be $r$-sample if
— Every pair of adjacent boundary samples p, q lies at a distance of at most $2 r$.
— Every pair of samples $p, q$ from the interior of $O$ lies at a minimum distance of $2 r$.

(a). r-sampling
(b). Point samples


## Sampling Models: Directed Boundary Sample

- Directed boundary sample is an r-sampling of object $O$ which possess a divergent boundary
- Theoretical analysis and topological correctness of RGG is provided under directed boundary sample
- Lemma: Let $S$ be a $(r, \uparrow)$-sample of an object $O, \partial R G G(S)$ contains an edge between every pair of adjacent samples of $\partial 0$.


## Sampling Models: Minimal Reach Sampling

- Interval $I(p)=\left[p_{0}, p_{1}\right]$ is the set of curve points between $p_{0}$, and $p_{1}$
- Reach of a curve interval I: inf Ifs(p):p $\in$



## Sampling Models: Minimal Reach Sampling ${ }^{[18]}$

- Consider pseudo-concavities with extremal bi-tangent points as the intervals
- Local feature size is computed w.r.t exterior medial axis
- Compute the minimum reach of $\gamma$ all the pseudo-concave intervals
- MRS: the closest neighboring point of any $p$ in $S$ lies at exactly
$\gamma$



Minimum Reach $=d_{3}=\min \left(d_{1}, d_{2}, d_{3}\right)$
[18] Thayyil et al. 2021, "A sampling type discernment approach towards reconstruction of a point set in $\mathrm{R}^{2 "}$, Computers Aided Geometric Design

## Evaluation Practices: $L^{2}$ error norm

- Quantitative analysis based on $L^{2}$ error norm ${ }^{[13]}$

$$
L^{2} \text { error norm }=\frac{\operatorname{area}((O-R e) \cup(R e-O))}{\operatorname{area}(O)}
$$

O: original object, Re: reconstructed polygon

- $L^{2}$ error norm of zero $\square$ both the areas are equal, and the boundaries are structurally alike.



## Evaluation Practices: Feature based Comparisons

- Typical features: simple closed curve, multiple components, holes, outliers, sharp corners



## Evaluation Practices: Point Set Density



Image Courtesy: [18]

## Evaluation Practices: Point Distributions

- DBDI: dense boundary and interior
- DBSI: sparse interior, dense bound
- SBDI: sparse boundary, dense int.
- SBSI: sparse boundary \& interior
- Other options: truly random, semi-random etc.
- Not robust to noise/outliers


Image Courtesy: [18]

## Chi Shape Software ${ }^{\left[{ }^{[3}\right]}$



- Point set generation of English letters and country maps


## Other Software

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- Alpha shape in CGAL Library
- C++ and CGAL predicates

| Sl. No | Algorithm | URL |
| :--- | :--- | :--- |
| 1 | CT-shape | https://github.com/agcl-mr/Reconstruction-CTShape |
| 2 | Petal ratio | https://github.com/agcl-mr/Reconstruction-Discern |
| 3 | Shape-hull graph | https://github.com/jijup/Shapehull2D |
| 4 | EC-shape | https://github.com/ShyamsTree/HoleDetection |

## Future Directions

- Improving and simplifying sampling conditions, especially for non-smooth and self-intersecting curves, and region reconstruction
- Reconstructing curves from hand drawn sketches with varying stroke thickness and intensity
- Deep learning on curves and shapes (similar to 2D medial axis)
- Reconstruct parametric curves instead of piece-wise polygonal curves
- Reconstruction of surfaces from networks of 3D curves
- Kinetic shapes: Shapes of moving points?


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#  

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